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# Introduction

**Dr. Nasir Jalal PhD (Computer Science)**

**Lecturer**

**CS & IT Department, CUVAS  
Bahawalpur**

# Course Detail



## BS CS 5<sup>th</sup> Semester Fall 2022-2026

Course: Theory of Programming Languages

Course Code: CS IT 507

## Lecture: 4 (18-09-2024)

# Outline Lecture 4



- Revision of Lecture 3
- Prerequisite of GFG
- What is Language
- Difference between Spoken language and Computer Language.
- What is CFG

# Symbols



Symbols are an entity or individual objects, which can be any letter, alphabet or any picture.

- Example:

1, a, b, #



# Alphabets

## Definition:

A finite non-empty set of symbols (letters), is called an alphabet. It is denoted by  $\Sigma$  ( Greek letter sigma).

## Example:

$$\Sigma = \{a, b\}$$

$\Sigma = \{0, 1\}$  //important as this is the language  
//which the computer understands.

$$\Sigma = \{i, j, k\}$$

- Examples:

$$\Sigma = \{a, b\}$$

$$\Sigma = \{A, B, C, D\}$$

$$\Sigma = \{0, 1, 2\}$$

$$\Sigma = \{0, 1, \dots, 5\}$$

$$\Sigma = \{\#, \beta, \Delta\}$$

# String

It is a finite collection of symbols from the alphabet. The string is denoted by  $w$ .

- **Example 1:**

If  $\Sigma = \{a, b\}$ , various string that can be generated from  $\Sigma$  are  $\{ab, aa, aaa, bb, bbb, ba, aba.....\}$ .

- **A string with zero occurrences of symbols is known as an empty string. It is represented by  $\epsilon$ .**
- **The number of symbols in a string  $w$  is called the length of a string. It is denoted by  $|w|$ .**

- **Example 2:**

$w = 010$

$|w| = 3$



# Words

- Definition:

Words are strings belonging to some language.

Example:

If  $\Sigma = \{f\}$  then a language  $L$  can be defined as  $L = \{f^n : n = 1, 2, 3, \dots\}$  or  $L = \{f, ff, fff, \dots\}$

Here  $f, ff, fff, \dots$  are the words of  $L$



# Valid/In-valid alphabets



- While defining an alphabet, an alphabet may contain letters consisting of group of symbols for example  $\Sigma_1 = \{B, aB, bab, d\}$ .
- Now consider an alphabet  $\Sigma_2 = \{B, Ba, bab, d\}$  and a string BababB.



# Scan “BababB”

This string can be **tokenized** in two different ways

- (Ba), (bab), (B)
- (B), (abab), (B)

Which shows that the second group cannot be identified as a string, defined over  $\Sigma$

# NOTE



- All words are strings, but not all strings are words.

# Language

A language is a collection of appropriate string. A language which is formed over  $\Sigma$  can be **Finite** or **Infinite**.

- **Example: 1**

$L1 = \{\text{Set of string of length 2}\} = \{aa, bb, ba, ab\}$

**Finite Language**

- **Example: 2**

$L2 = \{\text{Set of all strings starts with 'a'}\} = \{a, aa, aaa, abb, abbb, ababb\}$  **Infinite Language**



# Remarks

- While defining an alphabet of letters consisting of more than one symbols, no letter should be started with the letter of the same alphabet *i.e.* one letter should not be the prefix of another. However, a letter may be ended in the letter of same alphabet *i.e.* one letter may be the suffix of another.
- $\Sigma_1 = \{B, aB, bab, d\}$
- $\Sigma_2 = \{B, Ba, bab, d\}$

$\Sigma_1$  is a valid alphabet while  $\Sigma_2$  is an in-valid alphabet.

# Length of Strings

- Definition:

The length of string  $s$ , denoted by  $|s|$ , is the number of letters in the string.

- Example:

$\Sigma = \{a, b\}$

$s = ababa$

$|s| = 5$

# Example:2 String Length



$\Sigma = \{B, aB, bab, d\}$

$s = BaBbabBd$

Tokenizing = (B), (aB), (bab), (d)

$|s| = 4$

# Reverse of a String



- **Definition:**

The reverse of a string  $s$  denoted by  $\text{Rev}(s)$  or  $s^r$ , is obtained by writing the letters of  $s$  in reverse order.

- **Example:**

If  $s=abc$  is a string defined over  $\Sigma=\{a,b,c\}$   
then  $\text{Rev}(s)$  or  $s^r = cba$



# Reverse of a String



- Example 2:

$\Sigma = \{B, aB, bab, d\}$

$s = BaBbabBd$

$Rev(s) = dBbabaBB$



# CFG Overview

- A context-free grammar (CFG) consisting of a finite set of grammar rules is a quadruple

**Terminals** Represented using small case letter like a, b, c etc

**Non-Terminals**: The symbols that must be replaced by other things are called non-terminals. These symbols are represented using a capital letter like A, B, C, etc

**Productions**: The grammatical rules are often called productions or Principals.

**Start Symbol**:



# CFG Overview

A right-regular grammar is a **context-free grammar** in which the right-hand side of every production rule has one of the following forms: the empty string; a string consisting of a single non-terminal symbol; or a string consisting of a single terminal symbol followed by a single non-terminal symbol.



# CFG Overview

## CFG Example

$\Sigma = \{a,b\}$

productions:

1.  $S \rightarrow XaaX$
2.  $X \rightarrow aX$
3.  $X \rightarrow bX$
4.  $X \rightarrow \epsilon$

This grammar defines the language expressed by Regular Expression  $(a+b)^*aa(a+b)^*$ .

# CFG Example

Some of the strings generated by CFG are

aa

aabaababaabb

aaaa

aabaabbbbaabbb

aaabbbbbbbaabaaa

aaabbbbaaba →

aaaab

aaaaaabaab

aaaababbaabb

aabbbbaaaaaabaaa

Process to generate string aaabbbbaaba in explain in table

Rule	Application	Result
Start → S	Start	S
S → XaaX	S	XaaX
X → aX	XaaX	aXaaX
X → aX	aXaaX	aaXaaX
X → aX	aaXaaX	aaaXaaX
X → bX	aaaXaaX	aaabXaaX
X → bX	aaabXaaX	aaabbXaaX
X → bX	aaabbXaaX	aaabbbXaaX
X → ε	aaabbbXaaX	aaabbbbaaX
X → bX	aaabbbbaaX	aaabbbbaabX
X → aX	aaabbbbaabX	aaabbbbaabaX
X → ε	aaabbbbaabaX	aaabbbbaaba

# THANKS